



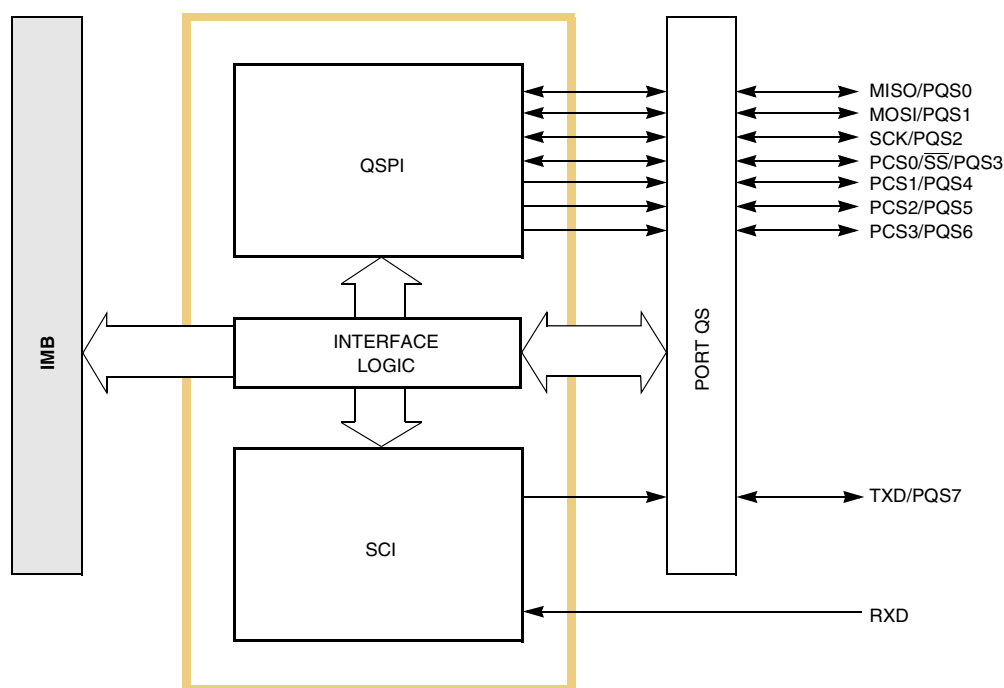
SECTION 6

QUEUED SERIAL MODULE

This section is an overview of the queued serial module (QSM). Refer to the [QSM Reference Manual](#) (QSMRM/AD) for complete information about the QSM. Refer to [APPENDIX D REGISTER SUMMARY](#) for a QSM address map and register bit and field definitions.

6.1 General

The QSM contains two serial interfaces, the queued serial peripheral interface (QSPI) and the serial communication interface (SCI). [Figure 6-1](#) is a block diagram of the QSM.



QSM BLOCK

Figure 6-1 QSM Block Diagram

The QSPI provides peripheral expansion or interprocessor communication through a full-duplex, synchronous, three-line bus. Four programmable peripheral chip-selects can select up to sixteen peripheral devices by using an external one of sixteen line selector. A self-contained RAM queue allows up to sixteen serial transfers of eight to sixteen bits each or continuous transmission of up to a 256-bit data stream without CPU32 intervention. A special wrap-around mode supports continuous transmission/reception modes.



The SCI provides a standard non-return to zero (NRZ) mark/space format. It operates in either full- or half-duplex mode. There are separate transmitter and receiver enable bits and dual data buffers. A modulus-type baud rate generator provides rates from 110 baud to 781 kbaud with a 25.17 MHz system clock. Word length of either eight or nine bits is software selectable. Optional parity generation and detection provide either even or odd parity check capability. Advanced error detection circuitry catches glitches of up to 1/16 of a bit time in duration. Wake-up functions allow the CPU32 to run uninterrupted until meaningful data is available.

6.2 QSM Registers and Address Map

There are four types of QSM registers: QSM global registers, QSM pin control registers, QSPI registers, and SCI registers. Refer to [6.2.1 QSM Global Registers](#) and [6.2.2 QSM Pin Control Registers](#) for a discussion of global and pin control registers. Refer to [6.3.1 QSPI Registers](#) and [6.4.1 SCI Registers](#) for further information about QSPI and SCI registers. Writes to unimplemented register bits have no effect, and reads of unimplemented bits always return zero. Refer to [5.2.1 Module Mapping](#) for more information about how the state of MM affects the system.

6.2.1 QSM Global Registers

The QSM configuration register (QSMCR) controls the interface between the QSM and the intermodule bus. The QSM test register (QTEST) is used during factory test of the QSM. The QSM interrupt level register (QILR) determines the priority of interrupts requested by the QSM and the vector used when an interrupt is acknowledged. The QSM interrupt vector register (QIVR) contains the interrupt vector for both QSM submodules. QILR and QIVR are 8-bit registers located at the same word address.

6.2.1.1 Low-Power Stop Mode Operation

When the STOP bit in QSMCR is set, the system clock input to the QSM is disabled and the module enters a low-power operating state. QSMCR is the only register guaranteed to be readable while STOP is asserted. The QSPI RAM is not readable during LPSTOP. However, writes to RAM or any register are guaranteed valid while STOP is asserted. STOP can be set by the CPU32 and by reset.

System software must bring the QSPI and SCI to an orderly stop before asserting STOP to avoid data corruption. The IRQ mask level in the CPU32 status register should be set to a higher value than the IRQ level generated by the QSM module. The SCI receiver and transmitter should be disabled after transfers in progress are complete. The QSPI can be halted by setting the HALT bit in SPCR3 and then setting STOP after the HALTA flag is set. The IRQ mask in the CPU status register should be restored to its former level. Refer to [5.3.4 Low-Power Operation](#) for more information about low-power stop mode.



6.2.1.2 Freeze Operation

The FRZ[1:0] bits in QSMCR are used to determine what action is taken by the QSM when the IMB FREEZE signal is asserted. FREEZE is asserted when the CPU32 enters background debug mode. At the present time, FRZ0 has no effect; setting FRZ1 causes the QSPI to halt on the first transfer boundary following FREEZE assertion. Refer to [4.10.2 Background Debug Mode](#) for more information.

6.2.1.3 QSM Interrupts

Both the QSPI and SCI can generate interrupt requests. Each has a separate interrupt request priority register. A single vector register is used to generate exception vector numbers.

The values of the ILQSPI and ILSCI fields in QILR determine the priority of QSPI and SCI interrupt requests. The values in these fields correspond to internal interrupt request signals $\overline{\text{IRQ}}[7:1]$. A value of %111 causes $\overline{\text{IRQ}}7$ to be asserted when a QSM interrupt request is made. Lower field values cause correspondingly lower-numbered interrupt request signals to be asserted. Setting the ILQSPI or ILSCI field values to %000 disables interrupts for the respective section. If ILQSPI and ILSCI have the same non-zero value, and the QSPI and SCI make simultaneous interrupt requests, the QSPI has priority.

When the CPU32 acknowledges an interrupt request, it places the value in the status register interrupt priority (IP) mask on the address bus. The QSM compares the IP mask value to the priority of the request to determine whether it should contend for arbitration priority. Arbitration priority is determined by the value of the IARB field in QSMCR. Each module that generates interrupts must have a non-zero IARB value. Arbitration is performed by means of serial contention between values stored in individual module IARB fields.

When the QSM wins interrupt arbitration, it responds to the CPU32 interrupt acknowledge cycle by placing an interrupt vector number on the data bus. The vector number is used to calculate displacement into the CPU32 exception vector table. SCI and QSPI vector numbers are generated from the value in the QIVR INTV field. The values of bits INTV[7:1] are the same for QSPI and SCI. The value of INTV0 is supplied by the QSM when an interrupt request is made. INTV0 = 0 for SCI interrupt requests; INTV0 = 1 for QSPI interrupt requests.

At reset, INTV[7:0] is initialized to \$0F, the uninitialized interrupt vector number. To enable interrupt-driven serial communication, a user-defined vector number must be written to QIVR, and interrupt handler routines must be located at the addresses pointed to by the corresponding vector. Writes to INTV0 have no effect. Reads of INTV0 return a value of one.



Refer to **SECTION 4 CENTRAL PROCESSOR UNIT** and **SECTION 5 SYSTEM INTEGRATION MODULE** for more information about exceptions and interrupts.

6.2.2 QSM Pin Control Registers

The QSM uses nine pins. Eight of the pins can be used for serial communication or for parallel I/O. Clearing a bit in the port QS pin assignment register (PQSPAR) assigns the corresponding pin to general-purpose I/O; setting a bit assigns the pin to the QSPI. PQSPAR does not affect operation of the SCI.

The port QS data direction register (DDRQS) determines whether pins are inputs or outputs. Clearing a bit makes the corresponding pin an input; setting a bit makes the pin an output. DDRQS affects both QSPI function and I/O function. DDQS7 determines the direction of the TXD pin only when the SCI transmitter is disabled. When the SCI transmitter is enabled, the TXD pin is an output.

The port QS data register (PORTQS) latches I/O data. PORTQS writes drive pins defined as outputs. PORTQS reads return data present on the pins. To avoid driving undefined data, first write PORTQS, then configure DDRQS.

PQSPAR and DDRQS are 8-bit registers located at the same word address. For a summary of QSM pin functions, refer to **Table 6-1**.



Table 6-1 Effect of DDRQS on QSM Pin Function

QSM Pin	QSPI Mode	DDRQS Bit	Bit State	Pin Function
MISO	Master	DDQS0	0	Serial data input to QSPI
			1	Disables data input
	Slave		0	Disables data output
			1	Serial data output from QSPI
MOSI	Master	DDQS1	0	Disables data output
			1	Serial data output from QSPI
	Slave		0	Serial data input to QSPI
			1	Disables data input
SCK ¹	Master	DDQS2	—	Clock output from QSPI
	Slave		—	Clock input to QSPI
PCS0/ \overline{SS}	Master	DDQS3	0	Assertion causes mode fault
			1	Chip-select output
	Slave		0	QSPI slave select input
			1	Disables slave select Input
PCS[1:3]	Master	DDQS[4:6]	0	Disables chip-select output
			1	Chip-select output
	Slave		0	Inactive
			1	Inactive
TXD ²	—	DDQS7	X	Serial data output from SCI
RXD	—	None	NA	Serial data input to SCI

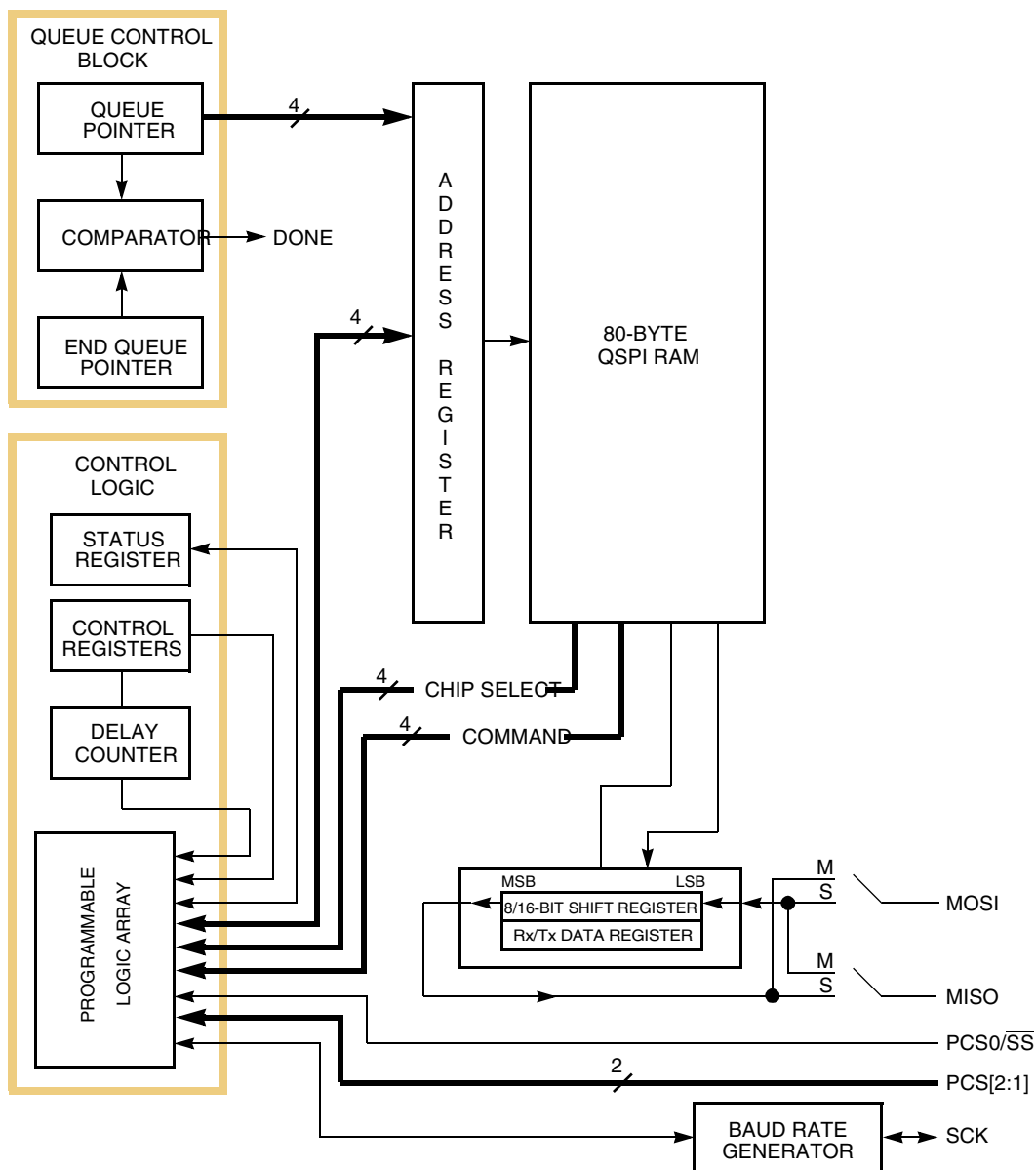
NOTES:

1. PQS2 is a digital I/O pin unless the SPI is enabled (SPE set in SPCR1), in which case it becomes the QSPI serial clock SCK.
2. PQS7 is a digital I/O pin unless the SCI transmitter is enabled (TE set in SCCR1), in which case it becomes the SCI serial data output TXD.

6.3 Queued Serial Peripheral Interface

The queued serial peripheral interface (QSPI) is used to communicate with external devices through a synchronous serial bus. The QSPI is fully compatible with SPI systems found on other Motorola products, but has enhanced capabilities. The QSPI can perform full duplex three-wire or half duplex two-wire transfers. A variety of transfer rates, clocking, and interrupt-driven communication options is available.

Figure 6-2 displays a block diagram of the QSPI.



QSPI BLOCK

Figure 6-2 QSPI Block Diagram

Serial transfers of eight to sixteen can be specified. Programmable transfer length simplifies interfacing to devices that require different data lengths.

An inter-transfer delay of 17 to 8192 system clocks can be specified (default is 17 system clocks). Programmable delay simplifies the interface to devices that require different delays between transfers.

A dedicated 80-byte RAM is used to store received data, data to be transmitted, and a queue of commands. The CPU32 can access these locations directly. This allows serial peripherals to be treated like memory-mapped parallel devices.



The command queue allows the QSPI to perform up to 16 serial transfers without CPU32 intervention. Each queue entry contains all the information needed by the QSPI to independently complete one serial transfer.

A pointer identifies the queue location containing the data and command for the next serial transfer. Normally, the pointer address is incremented after each serial transfer, but the CPU32 can change the pointer value at any time. Support of multiple-tasks can be provided by segmenting the queue.

The QSPI has four peripheral chip-select pins. The chip-select signals simplify interfacing by reducing CPU32 intervention. If the chip-select signals are externally decoded, 16 independent select signals can be generated.

Wrap-around mode allows continuous execution of queued commands. In wrap-around mode, newly received data replaces previously received data in the receive RAM. Wrap-around mode can simplify the interface with A/D converters by continuously updating conversion values stored in the RAM.

Continuous transfer mode allows transfer of an uninterrupted bit stream. Any number of bits in a range from 8 to 256 can be transferred without CPU32 intervention. Longer transfers are possible, but minimal intervention is required to prevent loss of data. A standard delay of 17 system clocks is inserted between the transfer of each queue entry.

6.3.1 QSPI Registers

The programmer's model for the QSPI consists of the QSM global and pin control registers, four QSPI control registers (SPCR[0:3]), the status register (SPSR), and the 80-byte QSPI RAM. Registers and RAM can be read and written by the CPU32. Refer to [APPENDIX D REGISTER SUMMARY](#) for register bit and field definitions.

6.3.1.1 Control Registers

Control registers contain parameters for configuring the QSPI and enabling various modes of operation. The CPU32 has read and write access to all control registers. The QSM has read access only to all bits except the SPE bit in SPCR1. Control registers must be initialized before the QSPI is enabled to ensure proper operation. SPCR1 must be written last because it contains the QSPI enable bit (SPE).

Writing a new value to any control register except SPCR2 while the QSPI is enabled disrupts operation. SPCR2 is buffered. New SPCR2 values become effective after completion of the current serial transfer. Rewriting NEWQP in SPCR2 causes execution to restart at the designated location. Reads of SPCR2 return the current value of the register, not of the buffer. Writing the same value into any control register except SPCR2 while the QSPI is enabled has no effect on QSPI operation.

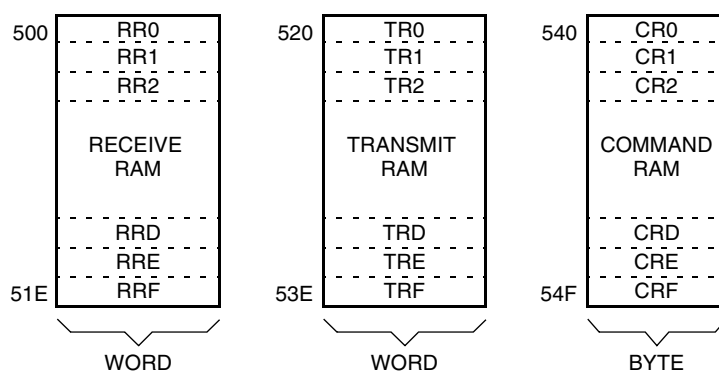
6.3.1.2 Status Register

SPSR contains information concerning the current serial transmission. Only the QSPI can set the bits in this register. The CPU32 reads SPSR to obtain QSPI status information and writes SPSR to clear status flags.



6.3.2 QSPI RAM

The QSPI contains an 80-byte block of dual-ported static RAM that can be accessed by both the QSPI and the CPU32. The RAM is divided into three segments: receive data RAM, transmit data RAM, and command data RAM. Receive data is information received from a serial device external to the MCU. Transmit data is information stored for transmission to an external device. Command control data defines transfer parameters. Refer to [Figure 6-3](#), which shows RAM organization.



QSPI RAM MAP

Figure 6-3 QSPI RAM

6.3.2.1 Receive RAM

Data received by the QSPI is stored in this segment. The CPU32 reads this segment to retrieve data from the QSPI. Data stored in the receive RAM is right-justified. Unused bits in a receive queue entry are set to zero by the QSPI upon completion of the individual queue entry. The CPU32 can access the data using byte, word, or long-word addressing.

The CPTQP value in SPSR shows which queue entries have been executed. The CPU32 uses this information to determine which locations in receive RAM contain valid data before reading them.



6.3.2.2 Transmit RAM

Data that is to be transmitted by the QSPI is stored in this segment and must be written by the CPU32 in right-justified form. The QSPI cannot modify information in the transmit RAM. The QSPI copies the information to its data serializer for transmission. Information remains in the transmit RAM until overwritten.

6.3.2.3 Command RAM

Command RAM is used by the QSPI in master mode. The CPU32 writes one byte of control information to this segment for each QSPI command to be executed. The QSPI cannot modify information in command RAM.

Command RAM consists of 16 bytes. Each byte is divided into two fields. The peripheral chip-select field enables peripherals for transfer. The command control field provides transfer options.

A maximum of 16 commands can be in the queue. Queue execution by the QSPI proceeds from the address in NEWQP through the address in ENDQP (both of these fields are in SPCR2).

6.3.3 QSPI Pins

The QSPI uses seven pins. These pins can be configured for general-purpose I/O when not needed for QSPI application.

Table 6-2 shows QSPI input and output pins and their functions.

Table 6-2 QSPI Pins

Pin Names	Mnemonics	Mode	Function
Master In Slave Out	MISO	Master Slave	Serial data input to QSPI Serial data output from QSPI
Master Out Slave In	MOSI	Master Slave	Serial data output from QSPI Serial data input to QSPI
Serial Clock	SCK	Master Slave	Clock output from QSPI Clock input to QSPI
Peripheral Chip Selects	PCS[3:1]	Master	Select peripherals
Slave Select	PCS0/ \overline{SS}	Master Master Slave	Selects peripherals Causes mode fault Initiates serial transfer

6.3.4 QSPI Operation

The QSPI uses a dedicated 80-byte block of static RAM accessible by both the QSPI and the CPU32 to perform queued operations. The RAM is divided into three segments. There are 16 command bytes, 16 transmit data words, and 16 receive data words. QSPI RAM is organized so that one byte of command data, one word of transmit data, and one word of receive data correspond to one queue entry, \$0-\$F.

The CPU32 initiates QSPI operation by setting up a queue of QSPI commands in command RAM, writing transmit data into transmit RAM, then enabling the QSPI. The QSPI executes the queued commands, sets a completion flag (SPIF), and then either interrupts the CPU32 or waits for intervention.



There are four queue pointers. The CPU32 can access three of them through fields in QSPI registers. The new queue pointer(NEWQP), contained in SPCR2, points to the first command in the queue. An internal queue pointer points to the command currently being executed. The completed queue pointer(CPTQP), contained in SPSR, points to the last command executed. The end queue pointer (ENDQP), contained in SPCR2, points to the final command in the queue.

The internal pointer is initialized to the same value as NEWQP. During normal operation, the command pointed to by the internal pointer is executed, the value in the internal pointer is copied into CPTQP, the internal pointer is incremented, and then the sequence repeats. Execution continues at the internal pointer address unless the NEWQP value is changed. After each command is executed, ENDQP and CPTQP are compared. When a match occurs, the SPIF flag is set and the QSPI stops and clears SPE, unless wrap-around mode is enabled.

At reset, NEWQP is initialized to \$0. When the QSPI is enabled, execution begins at queue address \$0 unless another value has been written into NEWQP. ENDQP is initialized to \$0 at reset, but should be changed to show the last queue entry before the QSPI is enabled. NEWQP and ENDQP can be written at any time. When NEWQP changes, the internal pointer value also changes. However, if NEWQP is written while a transfer is in progress, the transfer is completed normally. Leaving NEWQP and ENDQP set to \$0 transfers only the data in transmit RAM location \$0.

6.3.5 QSPI Operating Modes

The QSPI operates in either master or slave mode. Master mode is used when the MCU initiates data transfers. Slave mode is used when an external device initiates transfers. Switching between these modes is controlled by MSTR in SPCR0. Before entering either mode, appropriate QSM and QSPI registers must be initialized properly.

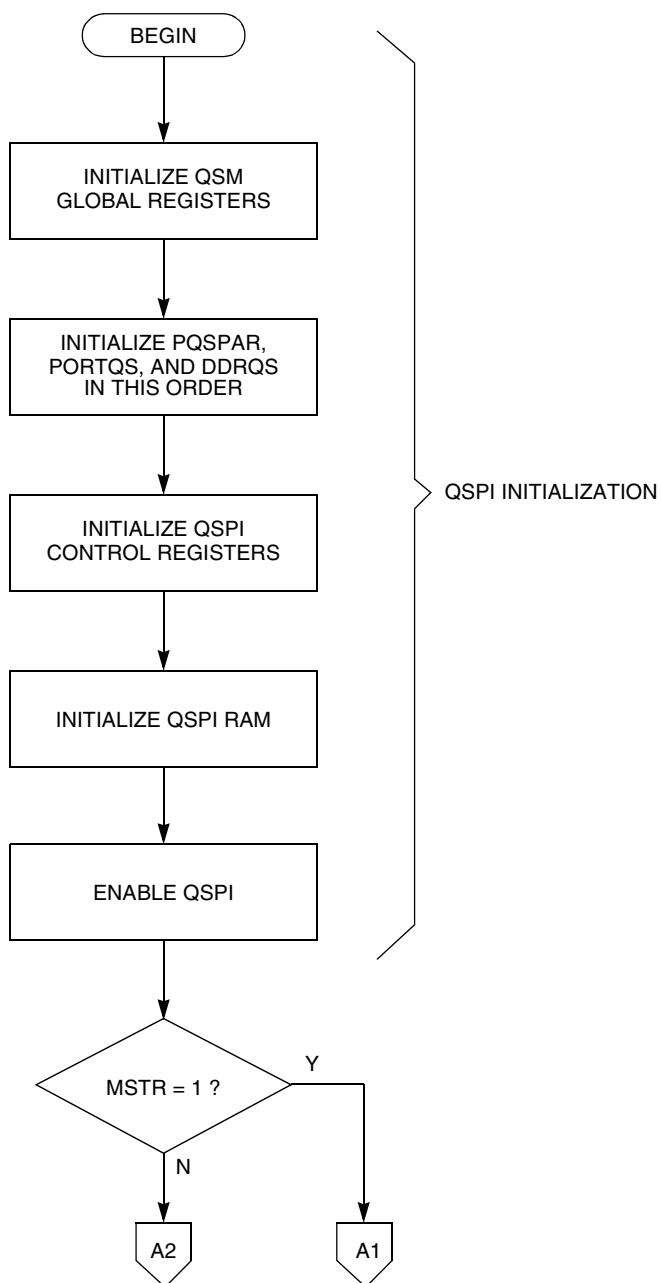
In master mode, the QSPI executes a queue of commands defined by control bits in each command RAM queue entry. Chip-select pins are activated, data is transmitted from the transmit RAM and received by the receive RAM.

In slave mode, operation proceeds in response to \overline{SS} pin activation by an external SPI bus master. Operation is similar to master mode, but no peripheral chip selects are generated, and the number of bits transferred is controlled in a different manner. When the QSPI is selected, it automatically executes the next queue transfer to exchange data with the external device correctly.

Although the QSPI inherently supports multi-master operation, no special arbitration mechanism is provided. A mode fault flag (MODF) indicates a request for SPI master arbitration. System software must provide arbitration. Note that unlike previous SPI systems, MSTR is not cleared by a mode fault being set nor are the QSPI pin output drivers disabled. The QSPI and associated output drivers must be disabled by clearing SPE in SPCR1.

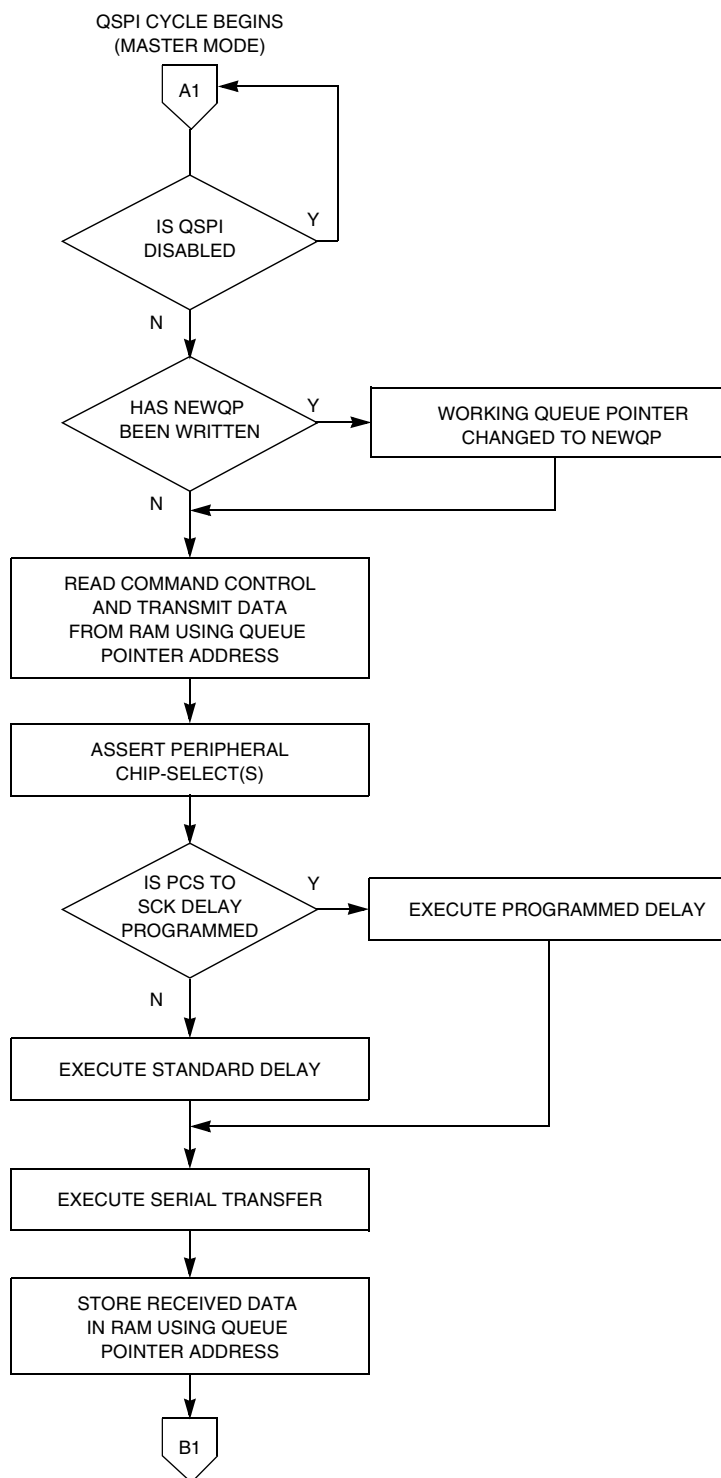


Figure 6-4 shows QSPI initialization. **Figure 6-5** through **Figure 6-9** show QSPI master and slave operation. The CPU32 must initialize the QSM global and pin registers and the QSPI control registers before enabling the QSPI for either mode of operation. The command queue must be written before the QSPI is enabled for master mode operation. Any data to be transmitted should be written into transmit RAM before the QSPI is enabled. During wrap-around operation, data for subsequent transmissions can be written at any time.



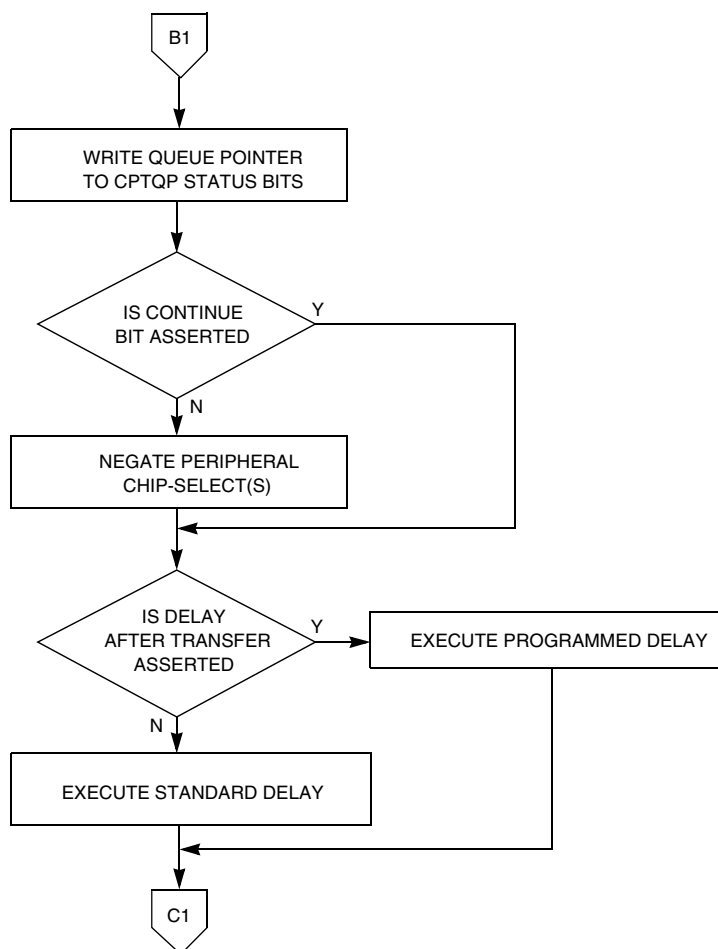
QSPI FLOW 1

Figure 6-4 Flowchart of QSPI Initialization Operation



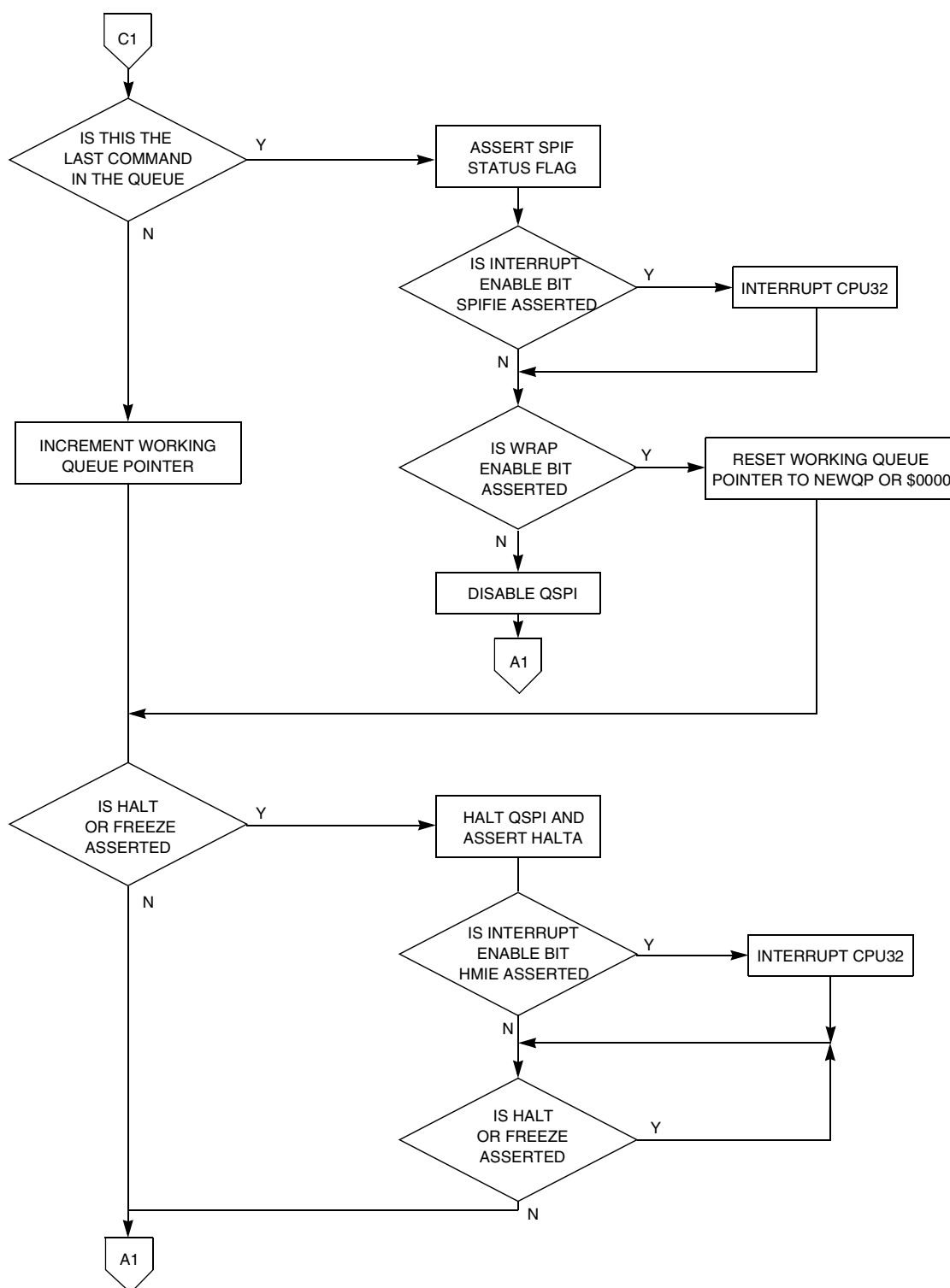
QSPI FLOW 2

Figure 6-5 Flowchart of QSPI Master Operation (Part 1)



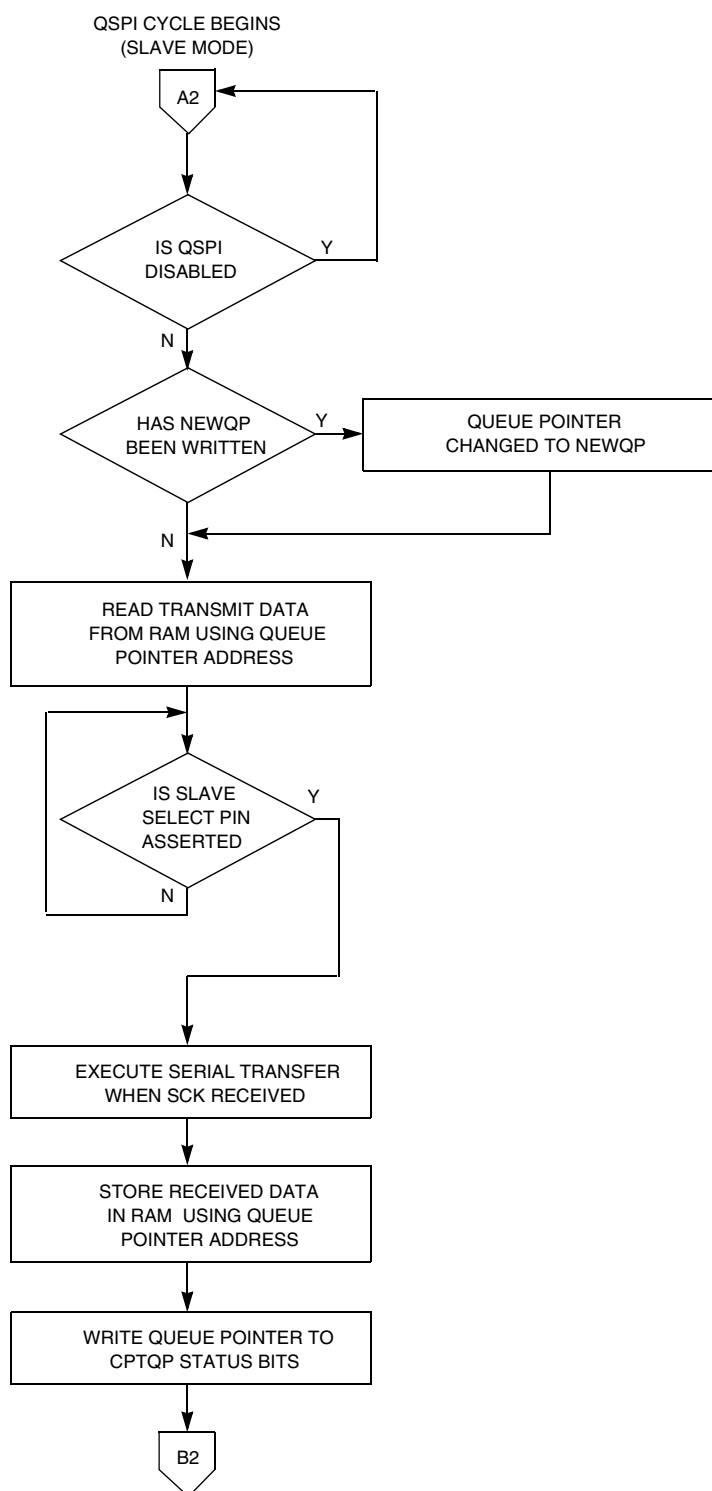
QSPI MSTR2 FLOW 3

Figure 6-6 Flowchart of QSPI Master Operation (Part 2)



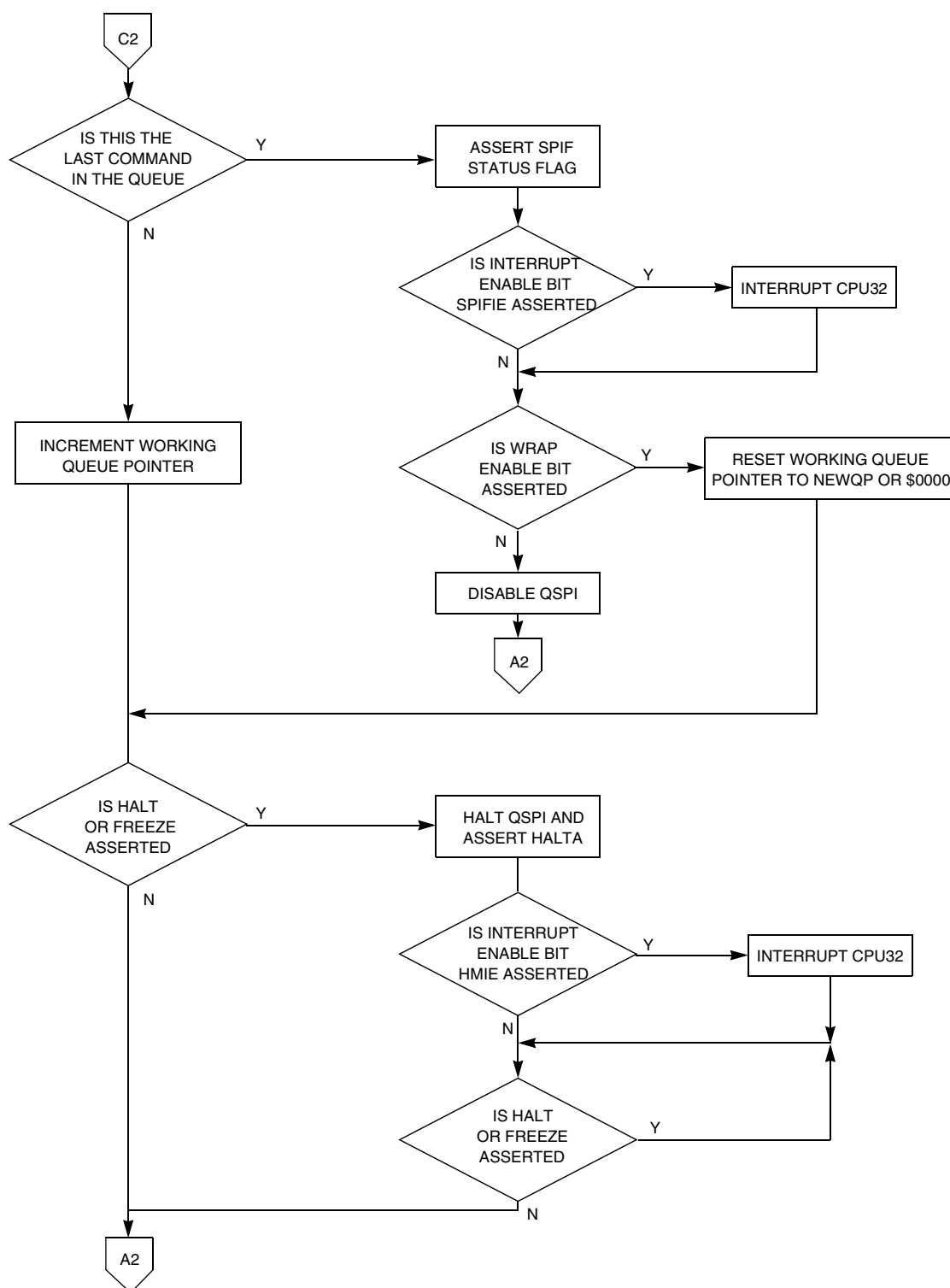
QSPI MSTR3 FLOW4

Figure 6-7 Flowchart of QSPI Master Operation (Part 3)



QSPI SLV1 FLOW 5

Figure 6-8 Flowchart of QSPI Slave Operation (Part 1)



QSPI SLV2 FLOW6

Figure 6-9 Flowchart of QSPI Slave Operation (Part 2)

Normally, the SPI bus performs synchronous bidirectional transfers. The serial clock on the SPI bus master supplies the clock signal SCK to time the transfer of data. Four possible combinations of clock phase and polarity can be specified by the CPHA and CPOL bits in SPCR0.



Data is transferred with the most significant bit first. The number of bits transferred per command defaults to eight, but can be set to any value from 8 to 16 bits inclusive by writing a value into the BITS[3:0] field in SPCR0 and setting BITSE in the command RAM.

Typically, SPI bus outputs are not open-drain unless multiple SPI masters are in the system. If needed, the WOMQ bit in SPCR0 can be set to provide wired-OR, open-drain outputs. An external pull-up resistor should be used on each output line. WOMQ affects all QSPI pins regardless of whether they are assigned to the QSPI or used as general-purpose I/O.

6.3.5.1 Master Mode

Setting the MSTR bit in SPCR0 selects master mode operation. In master mode, the QSPI can initiate serial transfers, but cannot respond to externally initiated transfers. When the slave select input of a device configured for master mode is asserted, a mode fault occurs.

Before QSPI operation begins, QSM register PQSPAR must be written to assign the necessary pins to the QSPI. The pins necessary for master mode operation are MISO, MOSI, SCK, and one or more of the chip-select pins. MISO is used for serial data input in master mode, and MOSI is used for serial data output. Either or both may be necessary, depending on the particular application. SCK is the serial clock output in master mode and must be assigned to the QSPI for proper operation.

The PORTQS data register must next be written with values that make the PQS2/SCK and PQS[6:3]/PCS[3:0] outputs inactive when the QSPI completes a series of transfers. Pins allocated to the QSPI by PQSPAR are controlled by PORTQS when the QSPI is inactive. PORTQS I/O pins driven to states opposite those of the inactive QSPI signals can generate glitches that momentarily enable or partially clock a slave device.

For example, if a slave device operates with an inactive SCK state of logic one (CPOL = 1) and uses active low peripheral chip-select PCS0, the PQS[3:2] bits in PORTQS must be set to %11. If PQS[3:2] = %00, falling edges will appear on PQS2/SCK and PQS3/PCS0 as the QSPI relinquishes control of these pins and PORTQS drives them to logic zero from the inactive SCK and PCS0 states of logic one.

Before master mode operation is initiated, QSM register DDRQS is written last to direct the data flow on the QSPI pins used. Configure the SCK, MOSI and appropriate chip-select pins PCS[3:0] as outputs. The MISO pin must be configured as an input.

After pins are assigned and configured, write appropriate data to the command queue. If data is to be transmitted, write the data to transmit RAM. Initialize the queue pointers as appropriate.

Data transfer is synchronized with the internally-generated serial clock SCK. Control bits, CPHA and CPOL, in SPCR0, control clock phase and polarity. Combinations of CPHA and CPOL determine upon which SCK edge to drive outgoing data from the MOSI pin and to latch incoming data from the MISO pin.



Baud rate is selected by writing a value from 2 to 255 into SPBR[7:0] in SPCR0. The QSPI uses a modulus counter to derive the SCK baud rate from the MCU system clock.

The following expressions apply to the SCK baud rate:

$$\text{SCK Baud Rate} = \frac{f_{\text{sys}}}{2 \times \text{SPBR}[7:0]}$$

or

$$\text{SPBR}[7:0] = \frac{f_{\text{sys}}}{2 \times \text{SCK Baud Rate Desired}}$$

Giving SPBR[7:0] a value of zero or one disables the baud rate generator and SCK assumes its inactive state.

The DSCK bit in each command RAM byte inserts either a standard (DSCK = 0) or user-specified (DSCK = 1) delay from chip-select assertion until the leading edge of the serial clock. The DSCKL field in SPCR1 determines the length of the user-defined delay before the assertion of SCK. The following expression determines the actual delay before SCK:

$$\text{PCS to SCK Delay} = \frac{\text{DSCKL}[6:0]}{f_{\text{sys}}}$$

where DSCKL[6:0] equals {1,2,3,..., 127}.

When DSCK equals zero, DSCKL[6:0] is not used. Instead, the PCS valid-to-SCK transition is one-half the SCK period.

There are two transfer length options. The user can choose a default value of eight bits, or a programmed value from 8 to 16 bits, inclusive. The programmed value must be written into BITS[3:0] in SPCR0. The BITSE bit in each command RAM byte determines whether the default value (BITSE = 0) or the BITS[3:0] value (BITSE = 1) is used. [Table 6-3](#) shows BITS[3:0] encoding.



Table 6-3 Bits Per Transfer

BITS[3:0]	Bits per Transfer
0000	16
0001	Reserved
0010	Reserved
0011	Reserved
0100	Reserved
0101	Reserved
0110	Reserved
0111	Reserved
1000	8
1001	9
1010	10
1011	11
1100	12
1101	13
1110	14
1111	15

Delay after transfer can be used to provide a peripheral deselect interval. A delay can also be inserted between consecutive transfers to allow serial A/D converters to complete conversion. Writing a value to DTL[7:0] in SPCR1 specifies a delay period. The DT bit in each command RAM byte determines whether the standard delay period (DT = 0) or the user-specified delay period (DT = 1) is used. The following expression is used to calculate the delay:

$$\text{Delay after Transfer} = \frac{32 \times \text{DTL}[7:0]}{f_{\text{sys}}} \text{ if DT} = 1$$

where DTL equals {1, 2, 3,..., 255}.

A zero value for DTL[7:0] causes a delay-after-transfer value of $8192/f_{\text{sys}}$.

$$\text{Standard Delay after Transfer} = \frac{17}{f_{\text{sys}}} \text{ if DT} = 0$$

Adequate delay between transfers must be specified for long data streams because the QSPI requires time to load a transmit RAM entry for transfer. Receiving devices need at least the standard delay between successive transfers. If the system clock is operating at a slower rate, the delay between transfers must be increased proportionately.

QSPI operation is initiated by setting the SPE bit in SPCR1. Shortly after SPE is set, the QSPI executes the command at the command RAM address pointed to by NEWQP. Data at the pointer address in transmit RAM is loaded into the data serializer and transmitted. Data that is simultaneously received is stored at the pointer address in receive RAM.



When the proper number of bits have been transferred, the QSPI stores the working queue pointer value in CPTQP, increments the working queue pointer, and loads the next data for transfer from transmit RAM. The command pointed to by the incremented working queue pointer is executed next, unless a new value has been written to NEWQP. If a new queue pointer value is written while a transfer is in progress, that transfer is completed normally.

When the CONT bit in a command RAM byte is set, PCS pins are continuously driven to specified states during and between transfers. If the chip-select pattern changes during or between transfers, the original pattern is driven until execution of the following transfer begins. When CONT is cleared, the data in register PORTQS is driven between transfers. The data in PORTQS must match the inactive states of SCK and any peripheral chip-selects used.

When the QSPI reaches the end of the queue, it sets the SPIF flag. SPIF is set during the final transfer before it is complete. If the SPIFIE bit in SPCR2 is set, an interrupt request is generated when SPIF is asserted. At this point, the QSPI clears SPE and stops unless wrap-around mode is enabled.

6.3.5.2 Master Wrap-Around Mode

Wrap-around mode is enabled by setting the WREN bit in SPCR2. The queue can wrap to pointer address \$0 or to the address pointed to by NEWQP, depending on the state of the WRTO bit in SPCR2.

In wrap-around mode, the QSPI cycles through the queue continuously, even while the QSPI is requesting interrupt service. SPE is not cleared when the last command in the queue is executed. New receive data overwrites previously received data in receive RAM. Each time the end of the queue is reached, the SPIF flag is set. SPIF is not automatically reset. If interrupt-driven QSPI service is used, the service routine must clear the SPIF bit to end the current interrupt request. Additional interrupt requests during servicing can be prevented by clearing SPIFIE, but SPIFIE is buffered. Clearing it does not end the current request.

Wrap-around mode is exited by clearing the WREN bit or by setting the HALT bit in SPCR3. Exiting wrap-around mode by clearing SPE is not recommended, as clearing SPE may abort a serial transfer in progress. The QSPI sets SPIF, clears SPE, and stops the first time it reaches the end of the queue after WREN is cleared. After HALT is set, the QSPI finishes the current transfer, then stops executing commands. After the QSPI stops, SPE can be cleared.

6.3.5.3 Slave Mode



Clearing the MSTR bit in SPCR0 selects slave mode operation. In slave mode, the QSPI is unable to initiate serial transfers. Transfers are initiated by an external SPI bus master. Slave mode is typically used on a multi-master SPI bus. Only one device can be bus master (operate in master mode) at any given time.

Before QSPI operation is initiated, QSM register PQSPAR must be written to assign necessary pins to the QSPI. The pins necessary for slave mode operation are MISO, MOSI, SCK, and PCS0/ \overline{SS} . MISO is used for serial data output in slave mode, and MOSI is used for serial data input. Either or both may be necessary, depending on the particular application. SCK is the serial clock input in slave mode and must be assigned to the QSPI for proper operation. Assertion of the active-low slave select signal (\overline{SS}) initiates slave mode operation.

Before slave mode operation is initiated, DDRQS must be written to direct data flow on the QSPI pins used. Configure the MOSI, SCK and PCS0/ \overline{SS} pins as inputs. The MISO pin must be configured as an output.

After pins are assigned and configured, write data to be transmitted into transmit RAM. Command RAM is not used in slave mode, and does not need to be initialized. Set the queue pointers, as appropriate.

When SPE is set and MSTR is clear, a low state on the slave select PCS0/ \overline{SS} pin begins slave mode operation at the address indicated by NEWQP. Data that is received is stored at the pointer address in receive RAM. Data is simultaneously loaded into the data serializer from the pointer address in transmit RAM and transmitted. Transfer is synchronized with the externally generated SCK. The CPHA and CPOL bits determine upon which SCK edge to latch incoming data from the MISO pin and to drive outgoing data from the MOSI pin.

Because the command RAM is not used in slave mode, the CONT, BITSE, DT, DSCK, and peripheral chip-select bits have no effect. The PCS0/ \overline{SS} pin is used only as an input.

The SPBR, DT and DSCKL fields in SPCR0 and SPCR1 bits are not used in slave mode. The QSPI drives neither the clock nor the chip-select pins and thus cannot control clock rate or transfer delay.

Because the BITSE option is not available in slave mode, the BITS field in SPCR0 specifies the number of bits to be transferred for all transfers in the queue. When the number of bits designated by BITS[3:0] has been transferred, the QSPI stores the working queue pointer value in CPTQP, increments the working queue pointer, and loads new transmit data from transmit RAM into the data serializer. The working queue pointer address is used the next time PCS0/ \overline{SS} is asserted, unless the CPU32 writes to NEWQP first.

The QSPI shifts one bit for each pulse of SCK until the slave select input goes high. If \overline{SS} goes high before the number of bits specified by the BITS field is transferred, the QSPI resumes operation at the same pointer address the next time \overline{SS} is asserted.

The maximum value that the BITS field can have is 16. If more than 16 bits are transmitted before \overline{SS} is negated, pointers are incremented and operation continues.



The QSPI transmits as many bits as it receives at each queue address, until the BITS[3:0] value is reached or \overline{SS} is negated. \overline{SS} does not need to go high between transfers as the QSPI transfers data until reaching the end of the queue, whether \overline{SS} remains low or is toggled between transfers.

When the QSPI reaches the end of the queue, it sets the SPIF flag. If the SPIFIE bit in SPCR2 is set, an interrupt request is generated when SPIF is asserted. At this point, the QSPI clears SPE and stops unless wrap-around mode is enabled.

6.3.5.4 Slave Wrap-Around Mode

Slave wrap-around mode is enabled by setting the WREN bit in SPCR2. The queue can wrap to pointer address \$0 or to the address pointed to by NEWQP, depending on the state of the WRTO bit in SPCR2. Slave wrap-around operation is identical to master wrap-around operation.

6.3.6 Peripheral Chip-Selects

Peripheral chip-select signals are used to select an external device for serial data transfer. Chip-select signals are asserted when a command in the queue is executed. Signals are asserted at a logic level corresponding to the value of the PCS[3:0] bits in each command byte. More than one chip-select signal can be asserted at a time, and more than one external device can be connected to each PCS pin, provided proper fanout is observed. PCS0 shares a pin with the slave select \overline{SS} signal, which initiates slave mode serial transfer. If \overline{SS} is taken low when the QSPI is in master mode, a mode fault occurs.

To configure a peripheral chip-select, set the appropriate bit in PQSPAR, then configure the chip-select pin as an output by setting the appropriate bit in DDRQS. The value of the bit in PORTQS that corresponds to the chip-select pin determines the base state of the chip-select signal. If base state is zero, chip-select assertion must be active high (PCS bit in command RAM must be set); if base state is one, assertion must be active low (PCS bit in command RAM must be cleared). PORTQS bits are cleared during reset. If no new data is written to PORTQS before pin assignment and configuration as an output, the base state of chip-select signals is zero and chip-select pins should thus be driven active-high.

6.4 Serial Communication Interface

The serial communication interface (SCI) communicates with external devices through an asynchronous serial bus. The SCI uses a standard non-return to zero (NRZ) transmission format. The SCI is fully compatible with other Motorola SCI systems, such as those on M68HC11 and M68HC05 devices. [Figure 6-10](#) is a block diagram of the SCI transmitter. [Figure 6-11](#) is a block diagram of the SCI receiver.

6.4.1 SCI Registers

The SCI programming model includes the QSM global and pin control registers, and four SCI registers. There are two SCI control registers (SCCR0 and SCCR1), one status register (SCSR), and one data register (SCDR). Refer to **APPENDIX D REGISTER SUMMARY** for register bit and field definitions.

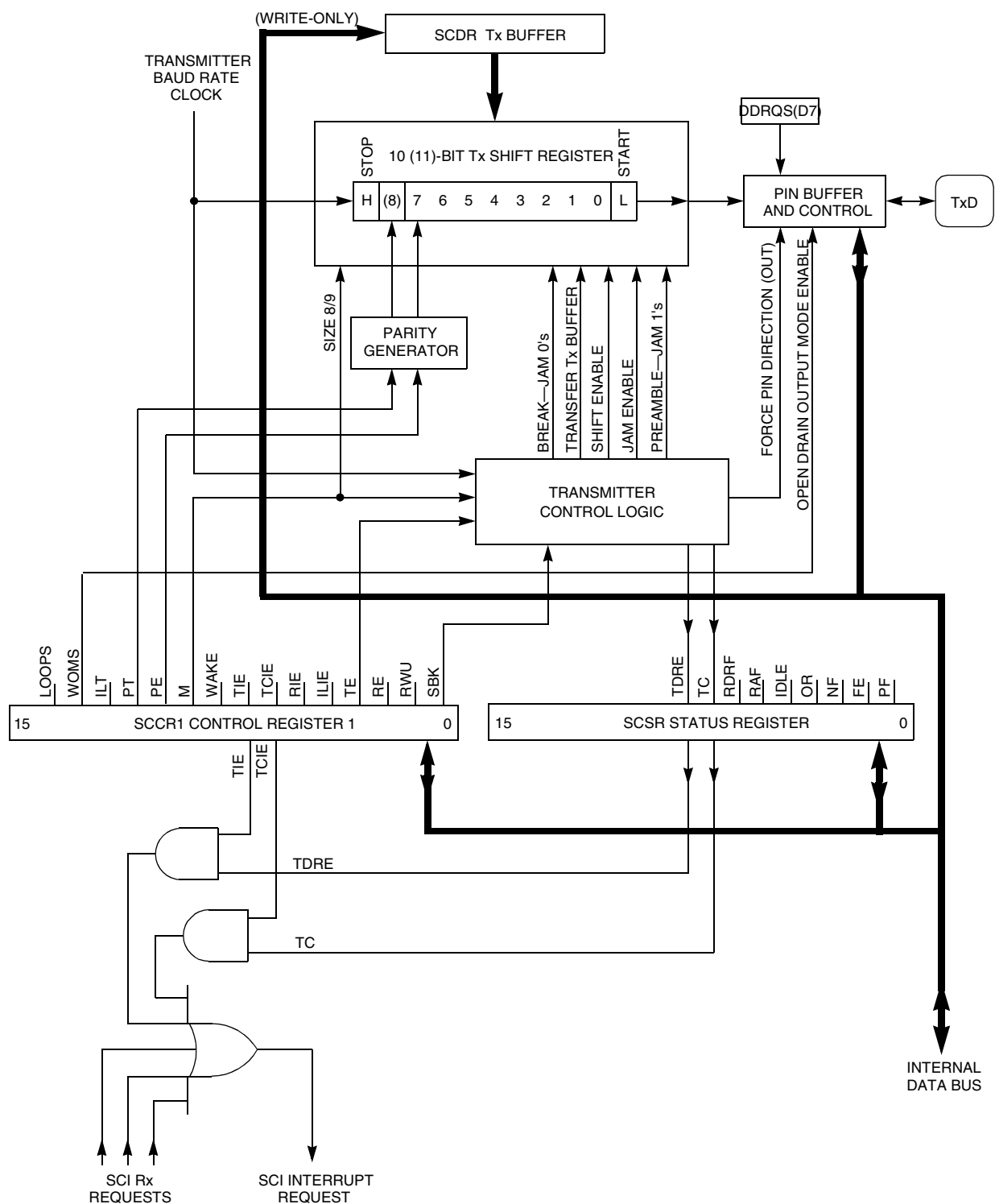


6.4.1.1 Control Registers

SCCR0 contains the baud rate selection field. Baud rate must be set before the SCI is enabled. This register can be read or written.

SCCR1 contains a number of SCI configuration parameters, including transmitter and receiver enable bits, interrupt enable bits, and operating mode enable bits. This register can be read or written at any time. The SCI can modify the RWU bit under certain circumstances.

Changing the value of SCI control bits during a transfer may disrupt operation. Before changing register values, allow the SCI to complete the current transfer, then disable the receiver and transmitter.



16/32 SCI TX BLOCK

Figure 6-10 SCI Transmitter Block Diagram

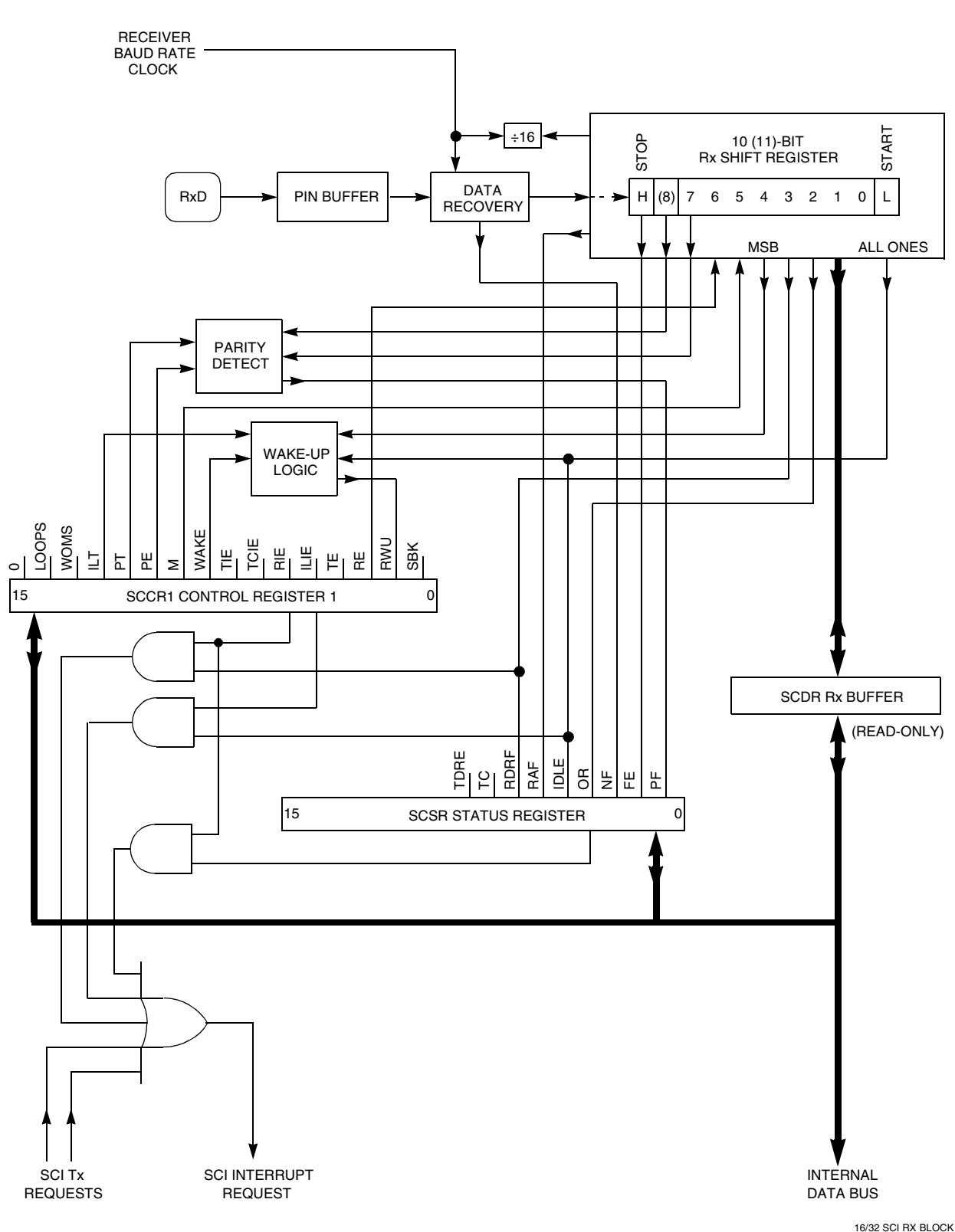


Figure 6-11 SCI Receiver Block Diagram

6.4.1.2 Status Register

SCSR contains flags that show SCI operating conditions. These flags are cleared either by SCI hardware or by reading SCSR, then reading or writing SCDR. A long-word read can consecutively access both SCSR and SCDR. This action clears receiver status flag bits that were set at the time of the read, but does not clear TDRE or TC flags.

If an internal SCI signal for setting a status bit comes after reading the asserted status bits, but before reading or writing SCDR, the newly set status bit is not cleared. SCSR must be read again with the bit set, and SCDR must be read or written before the status bit is cleared.

Reading either byte of SCSR causes all 16 bits to be accessed, and any status bit already set in either byte is cleared on a subsequent read or write of SCDR.

6.4.1.3 Data Register

SCDR contains two data registers at the same address. The receive data register (RDR) is a read-only register that contains data received by the SCI. Data enters the receive serial shifter and is transferred to RDR. The transmit data register (TDR) is a write-only register that contains data to be transmitted. Data is first written to TDR, then transferred to the transmit serial shifter, where additional format bits are added before transmission. R[7:0]/T[7:0] contain either the first eight data bits received when SCDR is read, or the first eight data bits to be transmitted when SCDR is written. R8/T8 are used when the SCI is configured for 9-bit operation. When the SCI is configured for 8-bit operation, they have no meaning or effect.

6.4.2 SCI Pins

Two unidirectional pins, TXD (transmit data) and RXD (receive data), are associated with the SCI. TXD can be used by the SCI or for general-purpose I/O. TXD function is controlled by PQSPA7 in the port QS pin assignment register (PQSPAR) and TE in SCI control register 1 (SCCR1). The receive data (RXD) pin is dedicated to the SCI.

6.4.3 SCI Operation

The SCI can operate in polled or interrupt-driven mode. Status flags in SCSR reflect SCI conditions regardless of the operating mode chosen. The TIE, TCIE, RIE, and ILIE bits in SCCR1 enable interrupts for the conditions indicated by the TDRE, TC, RDRF, and IDLE bits in SCSR, respectively.

6.4.3.1 Definition of Terms

- Bit-Time — The time required to transmit or receive one bit of data, which is equal to one cycle of the baud frequency.
- Start Bit — One bit-time of logic zero that indicates the beginning of a data frame. A start bit must begin with a one-to-zero transition and be preceded by at least three receive time samples of logic one.
- Stop Bit — One bit-time of logic one that indicates the end of a data frame.



- **Frame** — A complete unit of serial information. The SCI can use 10-bit or 11-bit frames.
- **Data Frame** — A start bit, a specified number of data or information bits, and at least one stop bit.
- **Idle Frame** — A frame that consists of consecutive ones. An idle frame has no start bit.
- **Break Frame** — A frame that consists of consecutive zeros. A break frame has no stop bits.



6.4.3.2 Serial Formats

All data frames must have a start bit and at least one stop bit. Receiving and transmitting devices must use the same data frame format. The SCI provides hardware support for both 10-bit and 11-bit frames. The M bit in SCCR1 specifies the number of bits per frame.

The most common data frame format for NRZ serial interfaces is one start bit, eight data bits (LSB first), and one stop bit; a total of 10 bits. The most common 11-bit data frame contains one start bit, eight data bits, a parity or control bit, and one stop bit. Ten-bit and eleven-bit frames are shown in [Table 6-4](#).

Table 6-4 Serial Frame Formats

10-Bit Frames			
Start	Data	Parity/Control	Stop
1	7	—	2
1	7	1	1
1	8	—	1
11-Bit Frames			
Start	Data	Parity/Control	Stop
1	7	1	2
1	8	1	1

6.4.3.3 Baud Clock

The SCI baud rate is programmed by writing a 13-bit value to the SCBR field in SCI control register zero (SCCR0). The baud rate is derived from the MCU system clock by a modulus counter. Writing a value of zero to SCBR[12:0] disables the baud rate generator. Baud rate is calculated as follows:

$$\text{SCI Baud Rate} = \frac{f_{\text{sys}}}{32 \times \text{SCBR}[12:0]}$$

or

$$\text{SCBR}[12:0] = \frac{f_{\text{sys}}}{32 \times \text{SCI Baud Rate Desired}}$$

where SCBR[12:0] is in the range {1, 2, 3, ..., 8191}.

The SCI receiver operates asynchronously. An internal clock is necessary to synchronize with an incoming data stream. The SCI baud rate generator produces a receive time sampling clock with a frequency 16 times that of the SCI baud rate. The SCI determines the position of bit boundaries from transitions within the received waveform, and adjusts sampling points to the proper positions within the bit period.



6.4.3.4 Parity Checking

The PT bit in SCCR1 selects either even (PT = 0) or odd (PT = 1) parity. PT affects received and transmitted data. The PE bit in SCCR1 determines whether parity checking is enabled (PE = 1) or disabled (PE = 0). When PE is set, the MSB of data in a frame is used for the parity function. For transmitted data, a parity bit is generated; for received data, the parity bit is checked. When parity checking is enabled, the PF bit in the SCI status register (SCSR) is set if a parity error is detected.

Enabling parity affects the number of data bits in a frame, which can in turn affect frame size. [Table 6-5](#) shows possible data and parity formats.

Table 6-5 Effect of Parity Checking on Data Size

M	PE	Result
0	0	8 data bits
0	1	7 data bits, 1 parity bit
1	0	9 data bits
1	1	8 data bits, 1 parity bit

6.4.3.5 Transmitter Operation

The transmitter consists of a serial shifter and a parallel data register (TDR) located in the SCI data register (SCDR). The serial shifter cannot be directly accessed by the CPU32. The transmitter is double-buffered, which means that data can be loaded into TDR while other data is shifted out. The TE bit in SCCR1 enables (TE = 1) and disables (TE = 0) the transmitter.

The shifter output is connected to the TXD pin while the transmitter is operating (TE = 1, or TE = 0 and transmission in progress). Wired-OR operation should be specified when more than one transmitter is used on the same SCI bus. The WOMS bit in SCCR1 determines whether TXD is an open-drain (wired-OR) output or a normal CMOS output. An external pull-up resistor on TXD is necessary for wired-OR operation. WOMS controls TXD function whether the pin is used by the SCI or as a general-purpose I/O pin.



Data to be transmitted is written to SCDR, then transferred to the serial shifter. The transmit data register empty (TDRE) flag in SCSR shows the status of TDR. When TDRE = 0, TDR contains data that has not been transferred to the shifter. Writing to SCDR again overwrites the data. TDRE is set when the data in TDR is transferred to the shifter. Before new data can be written to SCDR, however, the processor must clear TDRE by writing to SCSR. If new data is written to SCDR without first clearing TDRE, the data will not be transmitted.

The transmission complete (TC) flag in SCSR shows transmitter shifter state. When TC = 0, the shifter is busy. TC is set when all shifting operations are completed. TC is not automatically cleared. The processor must clear it by first reading SCSR while TC is set, then writing new data to SCDR.

The state of the serial shifter is checked when the TE bit is set. If TC = 1, an idle frame is transmitted as a preamble to the following data frame. If TC = 0, the current operation continues until the final bit in the frame is sent, then the preamble is transmitted. The TC bit is set at the end of preamble transmission.

The SBK bit in SCCR1 is used to insert break frames in a transmission. A non-zero integer number of break frames is transmitted while SBK is set. Break transmission begins when SBK is set, and ends with the transmission in progress at the time either SBK or TE is cleared. If SBK is set while a transmission is in progress, that transmission finishes normally before the break begins. To assure the minimum break time, toggle SBK quickly to one and back to zero. The TC bit is set at the end of break transmission. After break transmission, at least one bit-time of logic level one (mark idle) is transmitted to ensure that a subsequent start bit can be detected.

If TE remains set, after all pending idle, data and break frames are shifted out, TDRE and TC are set and TXD is held at logic level one (mark).

When TE is cleared, the transmitter is disabled after all pending idle data, and break frames are transmitted. The TC flag is set, and control of the TXD pin reverts to PQSPAR and DDRQS. Buffered data is not transmitted after TE is cleared. To avoid losing data in the buffer, do not clear TE until TDRE is set.

Some serial communication systems require a mark on the TXD pin even when the transmitter is disabled. Configure the TXD pin as an output, then write a one to PQS7. When the transmitter releases control of the TXD pin, it will revert to driving a logic one output.

To insert a delimiter between two messages, to place non-listening receivers in wake-up mode between transmissions, or to signal a retransmission by forcing an idle line, clear and then set TE before data in the serial shifter has shifted out. The transmitter finishes the transmission, then sends a preamble. After the preamble is transmitted, if TDRE is set, the transmitter will mark idle. Otherwise, normal transmission of the next sequence will begin.

Both TDRE and TC have associated interrupts. The interrupts are enabled by the transmit interrupt enable (TIE) and transmission complete interrupt enable (TCIE) bits in SCCR1. Service routines can load the last byte of data in a sequence into SCDR, then terminate the transmission when a TDRE interrupt occurs.



6.4.3.6 Receiver Operation

The RE bit in SCCR1 enables (RE = 1) and disables (RE = 0) the receiver. The receiver contains a receive serial shifter and a parallel receive data register (RDR) located in the SCI data register (SCDR). The serial shifter cannot be directly accessed by the CPU32. The receiver is double-buffered, allowing data to be held in RDR while other data is shifted in.

Receiver bit processor logic drives a state machine that determines the logic level for each bit-time. This state machine controls when the bit processor logic is to sample the RXD pin and also controls when data is to be passed to the receive serial shifter. A receive time clock is used to control sampling and synchronization. Data is shifted into the receive serial shifter according to the most recent synchronization of the receive time clock with the incoming data stream. From this point on, data movement is synchronized with the MCU system clock. Operation of the receiver state machine is detailed in the [QSM Reference Manual](#) (QSMRM/AD).

The number of bits shifted in by the receiver depends on the serial format. However, all frames must end with at least one stop bit. When the stop bit is received, the frame is considered to be complete, and the received data in the serial shifter is transferred to RDR. The receiver data register flag (RDRF) is set when the data is transferred.

Noise errors, parity errors, and framing errors can be detected while a data stream is being received. Although error conditions are detected as bits are received, the noise flag (NF), the parity flag (PF), and the framing error (FE) flag in SCSR are not set until data is transferred from the serial shifter to RDR.

RDRF must be cleared before the next transfer from the shifter can take place. If RDRF is set when the shifter is full, transfers are inhibited and the overrun error (OR) flag in SCSR is set. OR indicates that RDR needs to be serviced faster. When OR is set, the data in RDR is preserved, but the data in the serial shifter is lost. Because framing, noise, and parity errors are detected while data is in the serial shifter, FE, NF, and PF cannot occur at the same time as OR.

When the CPU32 reads SCSR and SCDR in sequence, it acquires status and data, and also clears the status flags. Reading SCSR acquires status and arms the clearing mechanism. Reading SCDR acquires data and clears SCSR.

When RIE in SCCR1 is set, an interrupt request is generated whenever RDRF or OR is set. Because receiver status flags are set at the same time as RDRF, they do not have separate interrupt enables.

6.4.3.7 Idle-Line Detection



During a typical serial transmission, frames are transmitted isochronally and no idle time occurs between frames. Even when all the data bits in a frame are logic ones, the start bit provides one logic zero bit-time during the frame. An idle line is a sequence of contiguous ones equal to the current frame size. Frame size is determined by the state of the M bit in SCCR1.

The SCI receiver has both short and long idle-line detection capability. Idle-line detection is always enabled. The idle line type (ILT) bit in SCCR1 determines which type of detection is used. When an idle line condition is detected, the IDLE flag in SCSR is set.

For short idle-line detection, the receiver bit processor counts contiguous logic one bit-times whenever they occur. Short detection provides the earliest possible recognition of an idle line condition, because the stop bit and contiguous logic ones before and after it are counted. For long idle-line detection, the receiver counts logic ones after the stop bit is received. Only a complete idle frame causes the IDLE flag to be set.

In some applications, software overhead can cause a bit-time of logic level one to occur between frames. This bit-time does not affect content, but if it occurs after a frame of ones when short detection is enabled, the receiver flags an idle line.

When the ILIE bit in SCCR1 is set, an interrupt request is generated when the IDLE flag is set. The flag is cleared by reading SCSR and SCDR in sequence. IDLE is not set again until after at least one frame has been received (RDRF = 1). This prevents an extended idle interval from causing more than one interrupt.

6.4.3.8 Receiver Wake-up

The receiver wake-up function allows a transmitting device to direct a transmission to a single receiver or to a group of receivers by sending an address frame at the start of a message. Hardware activates each receiver in a system under certain conditions. Resident software must process address information and enable or disable receiver operation.

A receiver is placed in wake-up mode by setting the RWU bit in SCCR1. While RWU is set, receiver status flags and interrupts are disabled. Although the CPU32 can clear RWU, it is normally cleared by hardware during wake-up.

The WAKE bit in SCCR1 determines which type of wake-up is used. When WAKE = 0, idle-line wake-up is selected. When WAKE = 1, address-mark wake-up is selected. Both types require a software-based device addressing and recognition scheme.

Idle-line wake-up allows a receiver to sleep until an idle line is detected. When an idle-line is detected, the receiver clears RWU and wakes up. The receiver waits for the first frame of the next transmission. The byte is received normally, transferred to RDR, and the RDRF flag is set. If software does not recognize the address, it can set RWU and put the receiver back to sleep. For idle-line wake-up to work, there must be a minimum of one frame of idle line between transmissions. There must be no idle time between frames within a transmission.

Address-mark wake-up uses a special frame format to wake up the receiver. When the MSB of an address-mark frame is set, that frame contains address information. The first frame of each transmission must be an address frame. When the MSB of a frame is set, the receiver clears RWU and wakes up. The byte is received normally, transferred to the RDR, and the RDRF flag is set. If software does not recognize the address, it can set RWU and put the receiver back to sleep. Address-mark wake-up allows idle time between frames and eliminates idle time between transmissions. However, there is a loss of efficiency because of an additional bit-time per frame.



6.4.3.9 Internal Loop Mode

The LOOPS bit in SCCR1 controls a feedback path in the data serial shifter. When LOOPS is set, the SCI transmitter output is fed back into the receive serial shifter. TXD is asserted (idle line). Both transmitter and receiver must be enabled before entering loop mode.

