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Initializing SDRAM  
Parameters for Motorola  
MPC106-Based Systems

This document describes the correlation of the programmable SDRAM interface parameters of Motorola's MPC106 PCI bridge/memory controller and typical SDRAM parameters found in manufacturers' data sheets. The MPC106 PCI bridge/memory controller provides a PowerPC™ microprocessor common hardware reference platform (CHRP™) compliant bridge between the PowerPC microprocessor family and the Peripheral Component Interconnect (PCI) bus.

Additional information about the topics discussed in this document can be found in *MPC106 PCI Bridge/Memory Controller User's Manual* (order #: MPC106UM/AD), and *Addendum to MPC106 PCI Bridge/Memory Controller User's Manual: MPC106 Revision 4.0 Supplement and User's Manual Errata* (order #: MPC106UMAD/AD). In this document, the term '106' is used as an abbreviation for, 'MPC106 PCI bridge/memory controller'.

To locate any published errata or updates for this document, refer to the website at <http://www.mot.com/SPS/PowerPC/>.

Note that this document describes the parameters for Rev. 4.0 of the MPC106. Earlier revisions are slightly different, but the information presented here is applicable with minor adjustment. For example, Rev. 4.0 has a 10-bit BSTOPRE parameter while Rev. 3.0 has an 8-bit BSTOPRE parameter. This means that while Rev. 4.0 devices can have a burst-to-precharge interval of 1023 clock cycles, Rev 3.0 devices are restricted to an interval of 255 clock cycles.

## 1 SDRAM Interface Parameters

At system reset, initialization software must set up the programmable parameters in the memory interface configuration registers (MICRs) of the 106. The MICRs control memory boundaries (starting and ending addresses), memory bank enables, memory buffer control, and memory controller operation and timing. Initialization software must program the MICRs at power-on reset and then enable the memory interface on the 106 by setting MCCR1[MEMGO].

The programmable parameters relevant to the SDRAM interface are:

- Memory bank starting and ending addresses (memory boundary registers)
- Memory bank enables (memory bank enable register)
- PGMAX—maximum activate-to-precharge interval
- SREN—self-refresh enable
- RAMTYP—RAM type

- PCKEN—memory interface parity checking/generation enable
- Bank *n* row—row address bit count for each bank
- BSTOPRE[0–9]—burst-to-precharge interval
- ECC\_EN—ECC enable
- REFINT—refresh interval
- BUF\_MODE—buffer mode
- RMW\_PAR—read-modify-write parity enable
- REFREC—refresh recovery interval
- RDLAT—data latency from read command
- PRETOACT—precharge-to-activate interval
- WCBUF—memory write buffer type
- RCBUF—memory read buffer type
- ACTOPRE—activate-to-precharge interval
- SDMODE—SDRAM mode register
- ACTORW—activate-to-read/write interval

The memory interface parameters, SREN, RAMTYP, PCKEN, ECC\_EN, BUF\_MODE, RMW\_PAR, WCBUF, and RCBUF are governed by the system design (refresh, parity, and buffering) and are not covered in this document.

## 2 Programming the Memory Boundary Registers

The extended starting address and the starting address registers are used to define the lower address boundary for each memory bank. The lower boundary is determined by the following formula:

Lower boundary for bank *n* = 0b00 || <extended starting address *n*> || <starting address *n*> || 0x0 0000

The extended ending address and the ending address registers are used to define the upper address boundary for each memory bank. The upper boundary is determined by the following formula:

Upper boundary for bank *n* = 0b00 || <extended ending address *n*> || <ending address *n*> || 0xF FFFF

Any unused banks should have their starting and ending addresses programmed out of the range of memory banks in use. If a disabled bank has its starting and ending address defined as overlapping an enabled bank's address space, there may be system memory corruption in the overlapping address range. Therefore, it is always important to map any unused bank's starting and ending address to memory space that is not used by the system.

For example, if a system has one 64-Mbyte bank of SDRAM, in bank 0, starting at address 0x0000\_0000 and ending at 0x03FF\_FFFF, the memory boundary registers should be programmed as shown in Table 1.

**Table 1. Memory Boundary Register Values for 64-Mbyte SDRAM in Bank 0**

Register	Value (Hex)	Address Offset (Hex)
Memory starting address 1	FFFF_FF00	80
Memory starting address 2	FFFF_FFFF	84

**Table 1. Memory Boundary Register Values for 64-Mbyte SDRAM in Bank 0 (continued)**

Register	Value (Hex)	Address Offset (Hex)
Extended memory starting address 1	0303_0300	88
Extended memory starting address 2	0303_0303	8C
Memory ending address 1	FFFF_FF3F	90
Memory ending address 2	FFFF_FFFF	94
Extended memory ending address 1	0303_0300	98
Extended memory ending address 2	0303_0303	9C

### 3 Programming the Memory Bank Enable Register

Individual banks are enabled or disabled by using the 1-byte memory bank enable register at address offset 0xA0. If a bank is enabled, the ending address of that bank must be greater than or equal to its starting address. As stated in the previous section, even if a bank is disabled, its starting and ending addresses must be programmed to memory space that is not used by enabled banks in the system.

### 4 Programming the Memory Page Mode Register

The 1-byte memory page mode register at address offset 0xA3 contains the PGMAX parameter, that controls how long the 106 retains the currently accessed page (row) in memory. The PGMAX parameter specifies the activate-to-precharge interval (sometimes called row active time or  $t_{RAS}$ ). The PGMAX value is multiplied by 64 to generate the actual number of clock cycles for the interval. When PGMAX is programmed to 0x00, page mode is disabled.

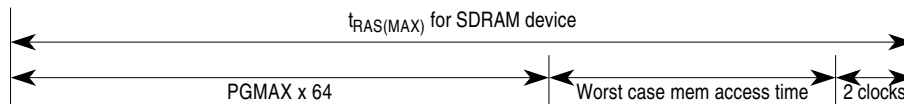
The value for PGMAX depends on the specific SDRAM devices used, the ROM system, and the operating frequency of the 106. When the interval specified by PGMAX expires, the 106 must close the active page by issuing a precharge bank command. PGMAX must be sufficiently less than the maximum row active time for the SDRAM device to ensure that the issuing of a precharge command is not stalled by a memory access. If a memory access is in progress when PGMAX expires, the 106 must wait for the access to complete before issuing the precharge command to the SDRAM. In the worst case, the 106 initiates a memory access one clock cycle before PGMAX expires. If ROM is located on the 60x/memory bus, the longest access that could potentially stall a precharge is a burst read from ROM. If ROM is located on the PCI bus, then it is not a factor, and the longest memory access that could potentially stall a precharge is a burst read from the SDRAM.

The 106 also requires two clock cycles to issue a precharge bank command to the SDRAM device. So, the PGMAX interval must be further reduced by two clock cycles.

Therefore, PGMAX should be programmed according to the following equation:

$$PGMAX < [t_{RAS(MAX)} - (\text{worst case memory access}) - 2] \div 64$$

PGMAX parameter settings are shown in Figure 1.


**Figure 1. PGMAX Parameter Setting for SDRAM Interface**

For example, consider a system with a 60x bus clock frequency of 66 MHz using SDRAMs with a maximum row active time ( $t_{RAS(MAX)}$ ) of 100  $\mu$ s. The maximum number of clock cycles between activate bank and precharge bank commands is 66 MHz  $\times$  100  $\mu$ s = 6600 clock cycles.

If the system uses 8-bit ROMs on the 60x/memory bus, a burst read from ROM follows the timing shown in Figure 6-40 in the *MPC106 PCI Bridge/Memory Controller User's Manual*. Also affecting the ROM access time is MCCR2[TS\_WAIT\_TIMER]. The minimum time allowed for ROM devices to enter high impedance is two clock cycles. TS\_WAIT\_TIMER adds clocks (n-1) to the minimum disable time. This delay is enforced after all ROM accesses preventing any other memory access from starting. Therefore, a burst read from an 8-bit ROM takes:

$$\{[(ROMFAL + 2) \times 8 + 3] \times 4 + 5\} + [2 + (TS\_WAIT\_TIMER - 1)] \text{ clock cycles}$$

So, if MCCR1[ROMFAL] = 4 and MCCR2[TS\_WAIT\_TIMER] = 3, the interval for a 60x burst read from an 8-bit ROM takes:

$$\{[(4 + 2) \times 8 + 3] \times 4 + 5\} + [2 + (3 - 1)] = 209 + 4 = 213 \text{ clock cycles}$$

Plugging the values into the PGMAX equation above:

$$PGMAX < (6600 - 213 - 2) \div 64 = 99.7 \text{ clock cycles}$$

The value stored in PGMAX should be 0b0110 0011 (or 99 clock cycles).

## 5 Programming the Memory Control Configuration Registers

The four 32-bit memory control configuration registers (MCCRs) set all RAM and ROM parameters. These registers are programmed by initialization software to adapt the 106 to the specific memory organization used in the system. The following sections describe programming individual SDRAM interface parameters within the MCCRs. Refer to the MCCR descriptions in Section 1.3.4, "Memory Control Configuration Registers," in the *Addendum to MPC106 PCI Bridge/Memory Controller User's Manual* and Section 3.2.6.4, "Memory Control Configuration Registers," in the *MPC106 PCI Bridge/Memory Controller User's Manual* for parameter locations within the MCCRs.

### 5.1 Bank *n* Row

Bits 15-0 of MCCR1 determine the row address configuration for the SDRAM devices in banks 0-7 (two bits per bank). For 64-Mbit devices with four internal banks, MCCR1[Bank *n* row] should be 0b00. Banks composed of 64-Mbit devices with two internal banks should have MCCR1[Bank *n* row] = 0b01. Banks composed of 16-Mbit devices with two internal banks should have MCCR1[Bank *n* row] = 0b11.

### 5.2 BSTOPRE[0-9]

The BSTOPRE parameter is composed of BSTOPRE[0-1] (bits 21-20 of MCCR2), BSTOPRE[2-5] (bits 31-28 of MCCR3), and BSTOPRE[6-9] (bits 3-0 of MCCR4). BSTOPRE controls the burst-to-precharge interval. BSTOPRE is similar to PGMAX in that when it expires, the 106 must generate a precharge

command. See Section 4, “Programming the Memory Page Mode Register,” for more information on the PGM\_MAX parameter. However, BSTOPRE is a much shorter duration counter that gets reloaded every time a read or write command is issued to the SDRAM devices.

The BSTOPRE interval can be optimized for the particular 106-based system implementation. If memory accesses are typically to the same rows within an active page, then a longer BSTOPRE interval would improve performance. Alternately, if memory accesses are typically to several locations, spanning multiple pages, then a shorter duration for BSTOPRE is in order. This allows for a precharge to close the active page before a subsequent access activates another page. Note that when BSTOPRE is programmed to 0b00 0000 0000, page mode is disabled.

### 5.3 REFINT

The memory interface supplies  $\overline{\text{CAS}}$  before  $\overline{\text{RAS}}$  (CBR) refreshes to SDRAM at the refresh interval specified by MCCR2[REFINT]. When REFINT expires, the 106 issues a precharge and then a refresh command to the SDRAM devices. The value stored in REFINT should allow for a potential collision between memory accesses and refresh cycles. In the worst case, the refresh must wait the number of clock cycles required by the longest access. For example, if ROM is located on the 60x/memory bus and a ROM access is in progress at the time a refresh operation needs to be performed, the refresh must wait until the ROM access has completed. If ROM is located on the 60x/memory bus, the longest access that could potentially stall a refresh is a burst read from ROM. If ROM is located on the PCI bus, the longest memory access is a burst read from the SDRAM.

The 106 also has to wait for a precharge command (to close any open pages) before it can issue the refresh command. The 106 requires two clock cycles to issue a precharge to an internal bank; with two pages open simultaneously, this equates to four extra clock cycles that must be taken off the refresh interval. Finally, the 106 must wait for the PRETOACT interval to pass before issuing the refresh command.

Therefore, REFINT should be programmed according to the following equation:

$$\text{REFINT} < (\text{per row refresh interval}) - (\text{worst case memory access}) - (\text{PRETOACT}) - 4$$

Consider a typical SDRAM device with a refresh period of 32 ms for a 2K cycle. This means that it takes 32 ms to refresh each internal bank and each internal bank has 2K rows. To refresh the whole SDRAM (two internal banks, 4K rows) it takes 64 ms. The refresh time per row is  $32 \text{ ms} \div 2048 \text{ rows}$  (or  $64 \text{ ms} \div 4096 \text{ rows}$ ) = 15.6  $\mu\text{s}$ . If the 60x bus clock is running at 66 MHz, the number of clock cycles per row refresh is  $15.6 \mu\text{s} \times 66 \text{ MHz} = 1030$  clock cycles.

If the system uses 8-bit ROMs on the 60x/memory bus, a burst read from ROM will follow the timing shown in Figure 6-40 in the 106 user’s manual. Also affecting the ROM access time is MCCR2[TS\_WAIT\_TIMER]. The minimum time allowed for ROM devices to enter high impedance is two clock cycles. TS\_WAIT\_TIMER adds clocks (n-1) to the minimum disable time. This delay is enforced after all ROM accesses preventing any other memory access from starting. Therefore, a burst read from an 8-bit ROM will take:

$$\{[(\text{ROMFAL} + 2) \times 8 + 3] \times 4 + 5\} + [2 + (\text{TS\_WAIT\_TIMER} - 1)] \text{ clock cycles}$$

So, if  $MCCR1[ROMFAL] = 4$  and  $MCCR2[TS\_WAIT\_TIMER] = 3$ , the interval for a 60x burst read from an 8-bit ROM will take:

$$\{[(4 + 2) \times 8 + 3] \times 4 + 5\} + [2 + (3 - 1)] = 209 + 4 = 213 \text{ clock cycles}$$

Plugging the values into the REFINT equation above:

$$\text{REFINT} < 1030 - 213 - 2 - 4 = 811 \text{ clock cycles}$$

The value stored in REFINT should be 0b00 0011 0010 1010 (or 810 clock cycles).

## 5.4 REFREC

REFREC is the refresh recovery interval. This parameter is the number of clock cycles from a refresh command until a bank activate command is allowed. The interval between a refresh and a new command is  $T_{RC}$  (row cycle time). A typical SDRAM data sheet shows  $T_{RC(MIN)} = 96 \text{ ns}$ . To put this time in clock cycles, for a 66-MHz bus,  $96 \text{ ns} \times 66 \text{ MHz} = 6.3$  clock cycles. Rounding up to 7, the value stored to REFREC should be 0b0111.

## 5.5 RDLAT

RDLAT is the data latency from a read command. This parameter controls the number of clock cycles from a read command until the first data beat is valid on the 60x data bus. The read command is issued with the assertion of  $\overline{SDCAS}$ . Therefore, this parameter equates to the  $\overline{CAS}$  latency ( $T_{CAC}$ ) plus whatever latency is imposed by the data buffers. Note that data latency is programmable for read operations (RDLAT). For write operations, the first valid data beat coincides with the SDRAM write command.

A typical SDRAM data sheet says that  $\overline{CAS}$  latency is 30 ns. To put this time in clock cycles, for a 66-MHz bus,  $30 \text{ ns} \times 66 \text{ MHz} = 1.98$  clock cycles (rounding up, 2 clock cycles). If the system uses registered data buffers, there is a one clock delay between data valid on the memory bus and data valid on the 60x bus, so the value stored to RDLAT should be  $2 + 1 = 3$  clock cycles (0b0011).

## 5.6 PRETOACT

PRETOACT is the number of clock cycles from a precharge command until a bank activate command is allowed.

A typical SDRAM data sheet shows that  $T_{RP}$  is the minimum row precharge time. That is, the time after a precharge command is issued before the bank activate command can begin.  $T_{RP}$  for a typical SDRAM device is 30 ns. To put this time in clock cycles, for a 66-MHz bus,  $30 \text{ ns} \times 66 \text{ MHz} = 1.98$  clock cycles. Rounding up to 2, the value stored to PRETOACT should be 0b0010.

## 5.7 ACTOPRE

ACTOPRE is the number of clock cycles from a bank activate command until a precharge command is allowed.

In a typical SDRAM data sheet,  $T_{RAS(MIN)}$  is the specification for minimum row active time. Every bank activate command must satisfy the  $T_{RAS(MIN)}$  specification before a precharge command can be issued to that active bank.

For a typical SDRAM device,  $T_{RAS(MIN)}$  is 70 ns. To put this time in clock cycles, for a 66-MHz bus,  $70 \text{ ns} \times 66 \text{ MHz} = 4.62$  clock cycles. Rounding up to 5, the value stored to ACTOPRE should be 0b0101.

## 5.8 SDMODE

SDMODE specifies the SDRAM mode register data to be written to the SDRAM array during power-up configuration. The 106 has certain restrictions for the SDMODE bits. In fact, the only variable in the SDMODE fields is the  $\overline{\text{CAS}}$  latency parameter.

The opcode should be 0b0 0000 for normal operation.

In a typical SDRAM data sheet,  $T_{\text{CAC}}$  is the specification for  $\overline{\text{CAS}}$  latency.  $T_{\text{CAC}}$  is programmable on most SDRAM devices to one, two, or three clock cycles depending on operating frequency of the device. A typical SDRAM data sheet says that  $\overline{\text{CAS}}$  latency is 30 ns. To put this time in clock cycles, for a 66-MHz bus,  $30 \text{ ns} \times 66 \text{ MHz} = 1.98$  clock cycles. Rounding up to 2, the value for  $\overline{\text{CAS}}$  latency should be 0b0010.

Finally, the wrap type must be 0b0 for sequential type wrapping and the wrap length must be 0b010 for four (doubleword) beats for each access.

Therefore, the complete value stored to SDMODE should be 0b0000 0010 0010.

Note that for 64-Mbit SDRAMs, the SDRAM mode register data is actually a 14-bit field. On the 106, the two-most-significant bits are forced to 0 and concatenated to MCCR4[SDMODE].

## 5.9 ACTORW

ACTORW controls the number of clocks cycles from a bank activate command until a read or write command is allowed. It must be at least two clock cycles.

In a typical SDRAM data sheet,  $T_{\text{RCD}}$  is the specification for the  $\overline{\text{RAS}}$ -to- $\overline{\text{CAS}}$  delay time. For SDRAM command encodings,  $\overline{\text{RAS}}$  corresponds to the bank activate command and  $\overline{\text{CAS}}$  corresponds to the read or write command.

For a typical SDRAM device,  $T_{\text{RCD}} = 25 \text{ ns}$ . To put this time in clock cycles, for a 66-MHz bus,  $25 \text{ ns} \times 66 \text{ MHz} = 1.65$  clock cycles. Rounding up to 2, the value for ACTORW should be 0b0010.

## 6 Revision History

Table 2 shows the revision history of this document.

**Table 2. Revision History**

Revision Number	Changes
0.0	Initial release
0.1	Nontechnical reformatting

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